



Optimality criterion for  $\forall y \in \mathcal{Y}$  is  $y^*$  is  $y \in \mathcal{Y}$

This is the first-order optimality condition.

In our case the problem is:

$$\begin{aligned} & \min_{z \in \mathcal{Z}} f_0(z) \\ & \text{subject to } z = \Pi_c(x) \\ & \quad \text{per unit cost change} \\ & \quad \text{total cost change as one moves from } y^* \text{ to any } y \text{ is positive if } y^* \text{ is optimal} \end{aligned}$$

$$\nabla f_0(z) = (z - x)$$

$$\nabla f_0(z^* = \Pi_c(x)) = (\Pi_c(x) - x)$$

So, the optimality condition is  $\forall z \in \mathcal{Z} \nabla f_0(z^*)^T (z - z^*) = (\Pi_c(x) - x)^T (z - \Pi_c(x)) \geq 0$

$$(\Pi_c(x) - x) (z - \Pi_c(x))$$

similarly when we are taking projection of  $y$ :

$$\forall z \in \mathcal{Z} \nabla f_0(z^*)^T (z - z^*) \geq 0$$

$$z = \Pi_c(y)$$

$$(\Pi_c(x) - x)^T (\Pi_c(y) - \Pi_c(x)) \geq 0$$

$$(\Pi_c(y) - y)^T (\Pi_c(x) - \Pi_c(y)) \geq 0$$

$$= ((\Pi_c(y) - y) + x - \Pi_c(x))^T (\Pi_c(x) - \Pi_c(y))$$

$$= ((x - y) - (\Pi_c(x) - \Pi_c(y)))^T (\Pi_c(x) - \Pi_c(y))$$

$$= (x - y)^T (\Pi_c(x) - \Pi_c(y)) - \|\Pi_c(x) - \Pi_c(y)\|_2^2 \geq 0$$

$$\rightarrow \|\Pi_c(x) - \Pi_c(y)\|_2^2 \leq (\Pi_c(x) - \Pi_c(y))^T (x - y) \leq \|\Pi_c(x) - \Pi_c(y)\|_2 \|x - y\|_2$$

By Cauchy-Schwarz  
 $a^T b \leq \|a\| \|b\|$

$$\|\Pi_c(x) - \Pi_c(y)\|_2 \leq \|x - y\|_2$$

$\therefore \Pi_c(\cdot)$  is a nonexpansive operator

Similarly, overprojection operator  $B_C = 2\Pi_C - I$  on  $C \subseteq \mathbb{R}^{n_p}$  is nonexpansive.

# Caution: nonexpansive and contraction mapping are function, but monotone or strongly monotone operator can be nontrivial relations.

A monotone operator always split out vector of same dimension as the input argument.

\* Monotone operators: def: monotone operators definitions and related

$$\text{relation: } R \text{ is monotone} \Leftrightarrow \forall (x, y) \in R \quad (x-y)^T (x-y) \geq 0 \quad \text{if } g \text{ for } f(x) = P^x \quad (x-y)^T (x-y) = f(x) - f(y) \geq 0$$

(in the increasing sense)

def: monotone operator  
 $\text{say } x \succ y \text{ then } x \succ y$

# Intuition of monotone operators via scalar functions

For that case monotone function is if input argument increases, the function value increases, and vice versa one way of quantifying that is:

$$(f(x) - f(y))(x-y) \geq 0 \Leftrightarrow (x-y) \geq 0 \Rightarrow (f(x) - f(y)) \geq 0$$

Now for imposing that as a definition we want this to be true for all  $(x, y)$  pair, so we can give the definition:

$$\forall x \in \text{dom } f \Leftrightarrow \forall y \in \text{dom } f \quad (f(x) - f(y))(x-y) \geq 0 \quad \# \text{ (+) direction of growth, (-) direction of decrease interpretation}$$

Note that (+) meaning  $x \succ y$  (y is growing), (-) meaning  $x \prec y$  (y is decreasing).

Now let's extend this for vector valued case, as vectors are partially ordered we cannot just start with  $x \succ y \Rightarrow f(x) \succ f(y)$ , however look at the scalar case  $\forall x, y \in \mathbb{R} \quad (f(x) - f(y))(x-y) \geq 0$  and trying to extend it point wise for a vector case:  $\forall x, y \in \mathbb{R}^n \quad \sum_i (f(x)_i - f(y)_i)(x_i - y_i) = (f(x) - f(y))^T (x - y) \geq 0$

What might be an intuition of this? Well it means that in some element of the vector the monotone operator is results in a such manner monotonic behavior to balance any anti-monotonic behavior in other elements. For a 2D case,  $(f(x) - f(y))_1 (x_1 - y_1) + (f(x) - f(y))_2 (x_2 - y_2) \geq 0$

So if the first term gets negative for some  $x, y$ , then the other terms will counter it by getting strongly positive. Note that this means:  $\forall x, y \in \text{dom } f \Leftrightarrow \forall u, v \in \text{dom } f \quad (u-v)^T (x-y) \geq 0$

- this is monotone, if we modify the notation as follows:

$$\forall u, v \in \text{dom } f \quad ((x, u) \in F \Leftrightarrow \forall y, v \in \text{dom } f \quad (u-v)^T (x-y) \geq 0) \quad \text{maximal monotone}$$

but it can be shown that, trivially  $\forall x \in \text{dom } f \quad (f(x) - f(y))(x-y) \geq 0 \wedge x \notin \text{dom } f \Rightarrow f(x) - f(y) = 0$  i.e.,  $f$  is monotone

∴ So, in summary we have a function  $f(x)$  which is monotone but not maximal monotone

\* Strong monotonicity:

$$\text{F}[\text{strongly monotone}] \Leftrightarrow \exists m > 0 \quad \forall (x, y) \in \text{dom } F \quad (u-v)^T (x-y) \geq m \|x-y\|_2^2$$

$\therefore \text{so } x \succ y \Leftrightarrow (u-v)^T (x-y) > 0$

$\forall x, y \in \text{dom } F \quad (f(x) - f(y)) \leq L \|x-y\|_2$  By Cauchy-Schwarz,  $(f(x) - f(y))^T (x-y) \leq \|f(x) - f(y)\|_2 \|x-y\|_2 \leq L \|x-y\|_2^2$

Frequently monotone, Lipschitz with constant  $L \geq \frac{1}{m}$   $m \|x-y\|_2^2 \leq (f(x) - f(y))^T (x-y) \leq L \|x-y\|_2^2$

$\Rightarrow L = \frac{1}{m} \geq 1$   
 condition number of strongly monotone Lipschitz function  $F$ .

\* Basic properties of monotone operator:

• Sum and scalar multiple

• F-monotone, G-monotone  $\Rightarrow F+G$  is monotone [eq: sum of monotone is monotone]

• F-monotone, G-monotone  $\Rightarrow F+G$  is maximal monotone # Additivity of maximal monotone operator with fine print

Some interior points of one relation domain has to belong to the other relation domain

• F[ maximal ] monotone  $\Rightarrow K \geq 0 F[\text{maximal}] \text{ monotone}$  # Positive scalar multiplicatively of (maximal) monotone operator

• F[ strongly monotone with parameter  $m \geq 0$  ]

$\Rightarrow \forall x, y \in \text{dom } F \quad (f(x) - f(y)) \geq m \|x-y\|_2^2$

• F[ strongly monotone ]  $\Rightarrow K \geq 0 F[\text{strongly monotone with parameter } m \geq 0]$

• Inverse:  $\exists^{-1}$  which always exists

• F[ maximal ] monotone  $\Rightarrow F^{-1}[ \text{maximal } ] \text{ monotone}$  [eq: inverse of monotone = monotone]

Frequently monotone with parameter  $m \geq 0 \Rightarrow F^{-1}$  function with derivative (instantaneous slope)  $m \geq 0$

Proof:  $\forall x, y \in \text{dom } F \quad (f(x) - f(y)) \geq m \|x-y\|_2^2$

Cauchy-Schwarz

$m \|x-y\|_2^2 \leq (f(x) - f(y))^T (x-y) \leq \|f(x) - f(y)\|_2 \|x-y\|_2$

definition of strong monotonicity

If we then:

WOW! # What does that mean? It means that: inverse of a strongly monotone operator (which may be a nontrivial relation) is a function (as any relation with a Lipschitz constant is a function) which is freaking amazing.

$$\begin{aligned} & \# \text{Intuition of monotone operators via scalar functions} \\ & \text{For that case monotone function is if input argument increases, the function value increases, and vice versa one way of quantifying that is:} \\ & (f(x) - f(y))(x-y) \geq 0 \Leftrightarrow (x-y) \geq 0 \Rightarrow (f(x) - f(y)) \geq 0 \\ & \text{Now for imposing that as a definition we want this to be true for all } (x, y) \text{ pair, so we can give the definition:} \\ & \forall x \in \text{dom } f \Leftrightarrow \forall y \in \text{dom } f \quad (f(x) - f(y))(x-y) \geq 0 \quad \# \text{ (+) direction of growth, (-) direction of decrease interpretation} \\ & \text{Note that (+) meaning } x \succ y \text{ (y is growing), (-) meaning } x \prec y \text{ (y is decreasing).} \\ & \text{Now let's extend this for vector valued case, as vectors are partially ordered we cannot just start with } x \succ y \Rightarrow f(x) \succ f(y), \text{ however look at the scalar case } \forall x, y \in \mathbb{R} \quad (f(x) - f(y))(x-y) \geq 0 \text{ and trying to} \\ & \text{extend it point wise for a vector case: } \forall x, y \in \mathbb{R}^n \quad \sum_i (f(x)_i - f(y)_i)(x_i - y_i) = (f(x) - f(y))^T (x - y) \geq 0 \\ & \text{What might be an intuition of this? Well it means that in some element of the vector the monotone operator is results in a such manner monotonic behavior to balance any anti-monotonic behavior in other elements. For a 2D case, } (f(x) - f(y))_1 (x_1 - y_1) + (f(x) - f(y))_2 (x_2 - y_2) \geq 0 \\ & \text{So if the first term gets negative for some } x, y, \text{ then the other terms will counter it by getting strongly positive. Note that this means: } \forall x, y \in \text{dom } f \Leftrightarrow \forall u, v \in \text{dom } f \quad (u-v)^T (x-y) \geq 0 \\ & \text{- this is monotone, if we modify the notation as follows:} \\ & \forall u, v \in \text{dom } f \quad ((x, u) \in F \Leftrightarrow \forall y, v \in \text{dom } f \quad (u-v)^T (x-y) \geq 0) \quad \text{maximal monotone} \end{aligned}$$

for example

$$\begin{aligned} & \text{this function is: } f(x) = x^2, \text{ dom } f = [0, \infty) \text{, so if a function is not maximal monotone then} \\ & \text{now set } x = 3 \notin \text{dom } f, \quad \forall y \in \text{dom } f \quad (f(x) - f(y))(x-y) \geq 0 \quad x \notin \text{dom } f \\ & (f(x) - f(y))(1-y) = (9-y^2)(3-y) = (3+y)(3-y)(3-y) = (3+y)(3-y) \geq 0 \\ & \exists x \in \mathbb{R} \quad \forall y \in \text{dom } f \quad (f(x) - f(y))(x-y) \geq 0 \quad \text{so, } f \text{ is not maximal monotone} \end{aligned}$$

but it can be shown that, trivially  $\forall x \in \text{dom } f \quad (f(x) - f(y))(x-y) \geq 0 \wedge x \notin \text{dom } f \Rightarrow f(x) - f(y) = 0$  i.e.,  $f$  is monotone

∴ So, in summary we have a function  $f(x)$  which is monotone but not maximal monotone



$\forall x \in \text{int}(C) \quad N_C(x) = \{0\}$

$\forall x \in C \quad N_C(x) \neq \text{nontrivial}$  # The intuition behind this is that remember at the boundary of any function, the subdifferential might not exist. One strange function in this regard is the function of  $\mathbb{R}^n$  set.

Why  $N_C(x)$  matters? Because this is the subdifferential mapping of the convex indicator function  $\gamma_C(x) = \begin{cases} 0, & x \in C \\ \infty, & x \notin C \end{cases}$ , i.e.,  $N_C = \gamma_C$ .

Makes sense because remember, for any convex function the subdifferential always exists in the int(dom f). At the boundary the subdifferential may not exist, i.e.,  $\gamma_C(x) \in N_C(x)$  might not exist, as a result  $N_C(x|_{\partial C}) = \gamma_C(x|_{\partial C}) = \text{nontrivial}$

[eq: subdifferential is monotone mapping]

proof:  $\gamma_C(x) = \begin{cases} 0, & x \in C \\ \infty, & x \notin C \end{cases}$ , want to determine,  $N_C(x) = N_\gamma(x)$

By definition of subgradient,  $\exists_{y \in \text{dom}(\gamma_C)} \forall_{y \in \text{dom}(\gamma_C)} \gamma_C(y) \geq \gamma_C(x) + g_x^T(y-x) \Leftrightarrow \min_{y \in \text{dom}(\gamma_C)} \gamma_C(y) = \gamma_C(x) + g_x^T(y-x) \Leftrightarrow \forall_{y \in \text{dom}(\gamma_C)} \gamma_C(y) \geq \gamma_C(x) + g_x^T(y-x)$

+ case 1:  $x \in C$ , then  $\forall_{y \in \text{dom}(\gamma_C)} 0 \geq \gamma_C(y) - \gamma_C(x) + g_x^T(y-x) \Leftrightarrow \forall_{y \in \text{dom}(\gamma_C)} g_x^T(y-x) \leq 0$

$\therefore x \in C \Rightarrow \forall_{y \in \text{dom}(\gamma_C)} g_x^T(y-x) \leq 0 \Rightarrow N_C(x) = \{0\}$

+ case 2:  $x \notin C \Rightarrow \exists_{y \in \text{dom}(\gamma_C)} 0 > \gamma_C(y) - \gamma_C(x) + g_x^T(y-x) \Rightarrow \text{no finite } g_x \text{ can exist} \Rightarrow N_C(x) = \emptyset$

Not the most rigorous proof ...

\* saddle subdifferential:

$f: \mathbb{R}^n \times \mathbb{R}^m \rightarrow \mathbb{R}$  s.t.  $\{(\cdot, \cdot)\}$  is convex in  $\mathbb{R}^n$ , concave in  $\mathbb{R}^m$

$F(x, y) = \{\text{saddle subdifferential relation}\}$

$$= \left[ \begin{array}{c} \partial_x f(x, y) \\ \partial_y f(x, y) \end{array} \right] \quad \forall_{x, y} \quad \partial_x f(x, y) \neq \emptyset, \partial_y f(x, y) \neq \emptyset$$

\* Why this is called saddle subdifferential?

Because  $(x, y) \in F \Leftrightarrow (x, y) \in \text{dom} f(x, y) \Leftrightarrow \begin{bmatrix} \partial_x f(x, y) \\ \partial_y f(x, y) \end{bmatrix} \Leftrightarrow \begin{array}{l} 0 \in \partial_x f(x, y) \Leftrightarrow x \in \text{argmin}_{\text{dom} f(x, y)} f(x, y) \Leftrightarrow \forall_{\tilde{x} \in \text{dom} f(x, y)} f(x, y) \leq f(\tilde{x}, y) \\ 0 \in \partial_y f(x, y) \Leftrightarrow y \in \text{argmin}_{\text{dom} f(x, y)} f(x, y) \Leftrightarrow \forall_{\tilde{y} \in \text{dom} f(x, y)} f(x, y) \leq f(x, \tilde{y}) \end{array}$

$\Leftrightarrow f(x, y) \geq f(\tilde{x}, y) \geq f(x, \tilde{y}) \geq f(\tilde{x}, \tilde{y})$

$\therefore \text{Set of } F \text{ is the saddle points of } f$ .

Saddle subdifferential relation for CCP function  $f(x, y)$  is maximal

\* KKT operator: # Manuscript version

$$\begin{pmatrix} \nabla f(x) \\ \nabla_h(x) \\ \nabla_l(x) \\ h(x) \\ l(x) \end{pmatrix} \quad \begin{array}{l} \text{affine} \\ \text{if } \lambda \neq 0 \\ \rightarrow 0 \end{array}$$

$$L(x, \lambda, v) = \begin{cases} f(x) + \sum_{i=1}^m \lambda_i h_i(x) + v_i l_i(x), & \text{if } \lambda \neq 0 \\ -\infty, & \text{otherwise} \end{cases}$$

$$T(x, \lambda, v) = \begin{bmatrix} \nabla_x L(x, \lambda, v) \\ -f(x) + \nabla_h(x) \\ \lambda \\ \nabla_l(x) \\ h(x) \end{bmatrix} \quad \begin{array}{l} \text{so } \nabla_x L(x, \lambda, v)^T \geq 0 \text{ (vanishing gradient of Lagrangian)} \\ \nabla_h(x) \geq 0 \text{ (primal inequality feasibility)} \\ \nabla_l(x) = 0 \text{ (dual inequality feasibility)} \\ \lambda \geq 0 \text{ (dual variable feasibility)} \\ \text{ensures that dual feasibility condition holds} \end{array}$$

$T(x, \lambda, v)$  is special case of saddle subdifferential if  $f$  is monotone

$\exists (x^*, \lambda^*, v^*) \in T(x^*, \lambda^*, v^*) \Leftrightarrow (x^*, \lambda^*, v^*)$  solves the optimality problem

\* KKT operator (slide version)

$$\nabla f(x)$$

$$Ax = b$$

$$L(x, y) = f(x) + y^T(Ax - b)$$

KKT operator:  $F(x, y) = \begin{bmatrix} \nabla_x L(x, y) \\ -\nabla_y L(x, y) \end{bmatrix}$  # gives vanishing gradient of Lagrangian (one of the KKT conditions)

Note that this is an equality constrained optimization problem, so it will have only these two KKT conditions with RHS zero and the LHS given by the rows of the KKT operator. So if the optimal  $(x^*, y^*)$  is inputted through the KKT operator then we will have zero.

$$F(0, 0) = \begin{bmatrix} \nabla_x L(0, 0) \\ -\nabla_y L(0, 0) \end{bmatrix}$$

so if  $F(x^*, y^*) = 0 \Leftrightarrow (x^*, y^*)$ : optimal primal-dual pair

$$\Leftrightarrow (x^*, y^*) \in F$$

$$\Leftrightarrow (0, (A^T y^*)) \in F^{-1}$$

$$\Leftrightarrow (x^*, y^*) \in F^{-1}(0)$$

So optimal primal-dual pair will belong to the output set of the inverse KKT operator with 0 fed into it!

# Multiplier to residual mapping is very important as it has a connection with ADMM

\* Multiplier to residual mapping. def: multiplier to residual mapping

$$\begin{pmatrix} \nabla f(x) \\ A \\ b \end{pmatrix} \quad \text{Residual}$$

$L(x, y) = f(x) + y^T(Ax - b)$  more technically:  $F(x) = b - A^T \text{argmin}_x L(x, y) \Leftrightarrow F(x) = b - A^T \text{argmin}_x f(x, y)$

def:  $F(x) = b - A^T x$  if  $x^* \in \text{argmin}_x L(x, y)$ , because  $f(x, y) = f(x^*, y) + \frac{1}{2} \|x - x^*\|^2$  since  $x \in \text{argmin}_x f(x, y)$  strongly convex, so will have unique minimizer

This is a kind of multiplier to residual mapping because it takes the lagrange multiplier and outputs the residual that associated with the sub-optimality of  $x$

Alternative definition to "residual mapping operator":  $F(y) = b - A^T x^*(y)$  Has a monotone operator, so multiplier to residual mapping is a monotone operator

$F(y) = b - A^T x^*(y)$  # monotone

# proof:  $\{\partial_x f(y)\}$  is monotone

$\Rightarrow \{\partial_x f(y)\}^{-1}$  is monotone

$\Rightarrow \forall_{y_1, y_2} \partial_x^T f(y_1) \leq \partial_x^T f(y_2) \Leftrightarrow \text{monotone}$

$\Rightarrow F: \text{monotone} \rightarrow A^T F(Ax) : \text{monotone}$

$$\downarrow$$

addition of a constant vector to each element of a relation set will not change monotonicity of a relation; equivalent logic. b1 ⊕ is a monotone operator  
 $-A(\delta_{x+y})^{-1}(-A^T \oplus)$  is a monotone operator }  $\oplus (b1 \oplus + (-A(\delta_{x+y})^{-1}(-A^T \oplus)))$  is a monotone operator  
 [ e.g. sum of monotone is monotone ]

NOW WE SHOW:

$$* F(y) = b - A(\beta_f)^{-1}(-A^T y)$$

$$= \partial_y [b^T y + f^*(-A^T y)]$$

## Proof:

$$\partial_y [ \partial_y f(x) (-h)] = \underbrace{\partial_y \partial_y f(x)}_{\text{---}} \underbrace{(-h)}_{\text{---}}$$

$$\text{now } \partial_y f^*(y) = \partial_y \left[ \sup_x -f(x) + y^T x \right]$$

$$= \partial_y \left[ -f(x^*) + y^T x^* \right] = \underset{x}{\text{argmax}} -f(x) + y^T x$$

} + recall, if  $f = \sup_{K \in A} (f_K)$

now  $\partial_y f^*(y) = \partial_y \left[ \sup_x [ -f(x) + y^T x ] \right]$

$= \partial_y \left[ -f(x^*) + y^T x^* \right] = \underset{x}{\arg\max} -f(x) + y^T x$  /\* recall, if  $f = \sup_{x \in A} f_A$

$x^* = \underset{x}{\arg\max} \{ f(x) \}_{x \in A}$   $= -f(x^*) + y^T x^* : x^* = \underset{x}{\arg\max} -f(x) + y^T x$   $\Rightarrow f = \sup_{x \in A} f_A = f_B$

$x^* = \underset{x}{\arg\min} f(x) - y^T x$  find  $\eta \in \partial f_B \Rightarrow \eta \in \partial f^*$

/\*  $x^* = \underset{x}{\arg\min} f(x) - y^T x$

$\Leftrightarrow [\partial_x (f(x) - y^T x)]_{x=x^*} \ni 0$

$\Leftrightarrow \partial f(x^*) - y \ni 0 \Leftrightarrow \partial f(x^*) \ni y$

$\Leftrightarrow (x^*, y) \in \partial f \Leftrightarrow (y, x^*) \in (\partial f)^{-1} \Leftrightarrow x^* = (\partial f^{-1})(y)$  \*/

this is a function  
(inverse of subdifferential  
is a function)

$= (\partial f^{-1})(y)$

$$\therefore \partial_y (b^T y + f^*(-A^T y)) = b - \partial_y (f^*(-A^T y)) = b - [(2f^{-1})'(y)]_{y := -A^T y}$$

$$= b - A(\partial f^{-1})(-A^T y) = b - A[\partial f^{-1}(y)]_{y=-A^T y} = b - A(\partial f^{-1})(-A^T y) \quad \square$$